TCP

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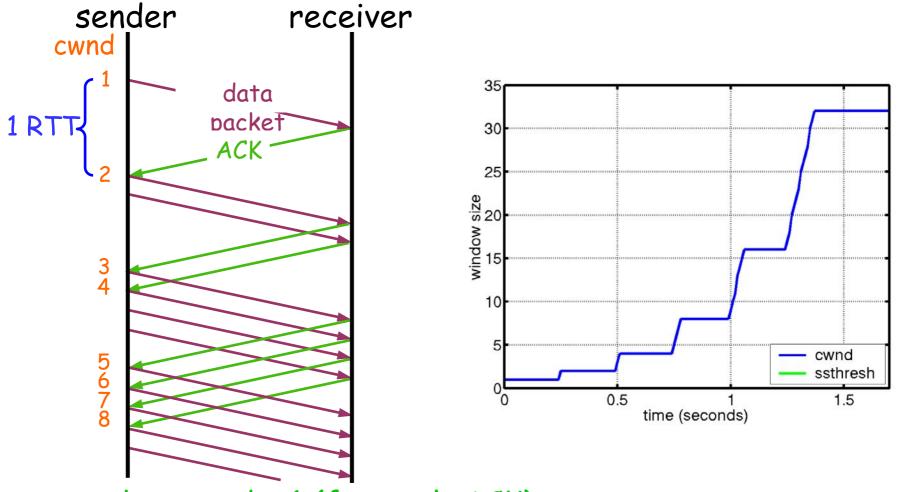
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TCP Congestion Control

Has four main parts Slow Start (SS) Congestion Avoidance (CA) ► Tahoe Fast Retransmit Fast Recovery ssthresh: slow start threshold determines whether to use SS or CA Assume packet losses are caused by congestion

Slow Start

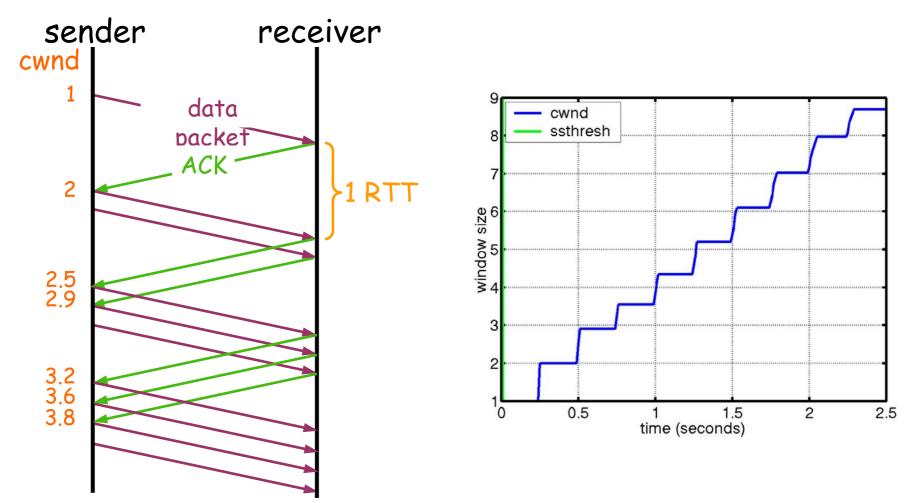


 $cwnd \leftarrow cwnd + 1$ (for each ACK)

Congestion Avoidance

 Starts when cwnd ≥ ssthresh
 On each successful ACK: cwnd ← cwnd + 1/cwnd
 Linear growth of cwnd each RTT: cwnd ← cwnd + 1

Congestion Avoidance



 $cwnd \leftarrow cwnd + 1/cwnd$ (for each ACK)

Packet Loss

Assumption: loss indicates congestion
 Packet loss detected by

 Retransmission TimeOuts (RTO timer)
 Duplicate ACKs (at least 3)

Packets

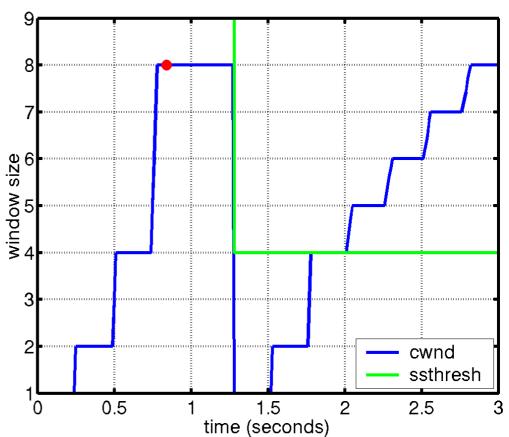
1 2 3 5 6 7

Acknowledgements

 1
 2
 3
 3
 3
 3

Timeout

ssthresh \leftarrow cwnd/2 cwnd = 1



Fast Retransmit

Wait for a timeout is quite long

- Immediately retransmits after 3 dupACKs without waiting for timeout
- Adjusts ssthresh

flightsize = min(awnd, cwnd)

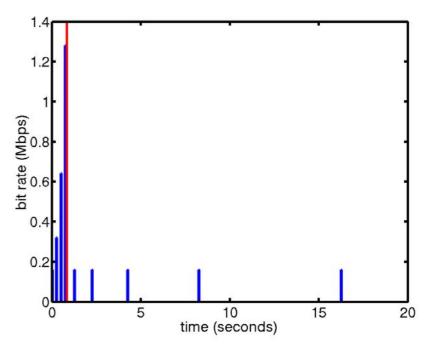
ssthresh \leftarrow max(flightsize/2, 2)

Enter Slow Start (cwnd = 1)

Successive Timeouts

When there is a timeout, double the RTO

- Keep doing so for each lost retransmission
 - Exponential back-off
 - Max 64 seconds¹
 - Max 12 restransmits¹



1 - Net/3 BSD

Fast recovery

- Motivation: prevent `pipe' from emptying after fast retransmit
- Idea: each dupACK represents a packet having left the pipe (successfully received)
- Enter FR/FR after 3 dupACKs
 - Set ssthresh \leftarrow max(flightsize/2, 2)
 - Retransmit lost packet

 - Wait till W=min(awnd, cwnd) is large enough; transmit new packet(s)

Enter CA

TCP Reno

